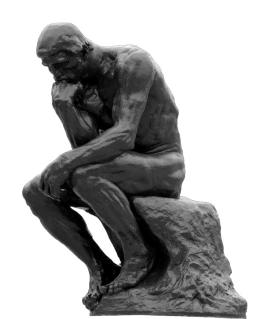
CS 161 Design and Analysis of Algorithms

Lecture 1:

Logistics, introduction, and multiplication!

The big questions

- Who are we?
 - Course staff, students?
- Why are we here?
 - Why learn about algorithms?
- What is going on?
 - What is this course about?
 - Logistics?
- Can we multiply integers?
 - And can we do it quickly?



Who are we?

- Instructor:
 - **Mary Wootters**
- **Course Coodinator:**
 - Amelie Byun
- **Embedded EthiCS Team:**
 - Justin Shin, Jennifer Chien, **Louis Ortiz**



- Anisha Palaparthi (Head CA)
- Maya Avital (Embedded EthiCS CA)
- Bradley Moon (Student Liaison)
- Zayn Malhotra
- Ta-Wei Tu
- Mingwei Yang
- **Spencer Compton**
- Will Fang
- Ziyi Ding

Ly-Ly Atchariyachanvanit

Justin

Mary

Amelie

- Karan Bhasin
- Simon Kim
- Ziyi Ding
 - James Cheng







Anisha

Isabel

Karan

Ly-Ly







Ta-Wei

Jennifer

Spencer

Simon







Will

Zayn

Maya









Xiao

Mingwei

James





Bradley

Andy

Auddithio

Andy Dai

Louie

Isabel Sieh

Xiao Mao

Auddithio Nag

Who are you?

- Freshman
- Sophomores
- Juniors
- Seniors

 MA/MS Students

- PhD Students
- NDO Students

Concentrating in:

- Art Practice
- Bioengineering
- Biology
- Biomedical Data Science
- Biomedical Informatics
- Chemical Eng.
- Chemistry
- Civil & Env. Eng.

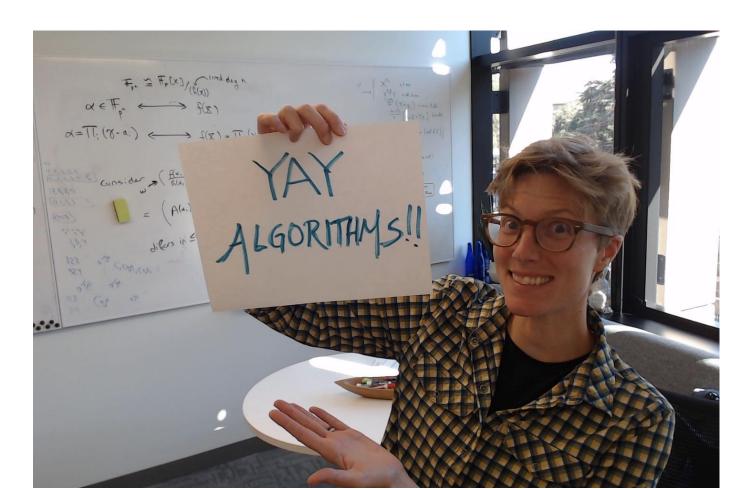
- Classics
- Communication
- CME
- Computer Science
- Creative Writing
- Data Science
- Earth Systems
- Economics
- Education

- EE
- Engineering
- Hum Bio
- International Relations
- Linguistics
- Math
- Music
- MS&E
- Mech. Eng.

- Physics
- Political Science
- Sociology
- Statistics
- Symbolic Systems
- Theater and Perf. Studies
- Undeclared

Why are we here?

• I'm here because I'm super excited about algorithms!



You are better equipped to answer this question than I am, but I'll give it a go anyway...

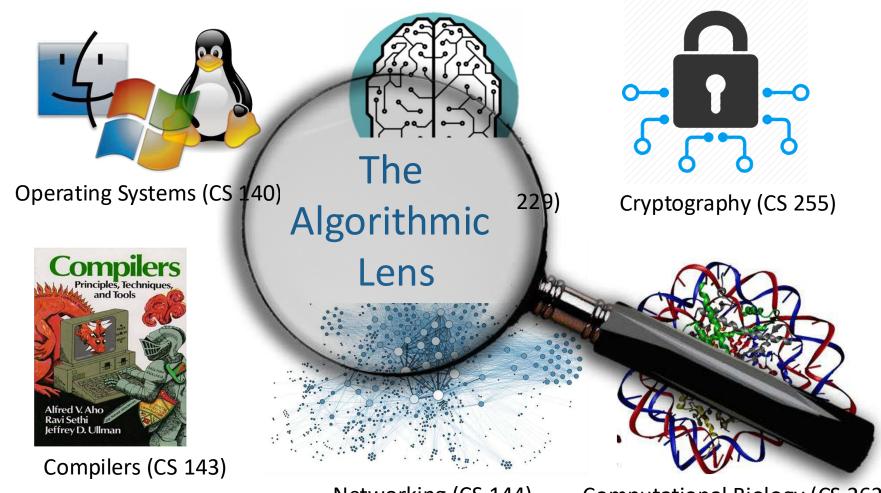
Why are you here?

- Algorithms are fundamental.
- Algorithms are useful.
- Algorithms are fun!
- CS161 is a required course.

Why is CS161 required?

- Algorithms are fundamental.
- Algorithms are useful.
- Algorithms are fun!

Algorithms are fundamental



Networking (CS 144)

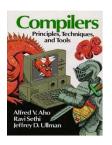
Computational Biology (CS 262)

Algorithms are useful

- All those things without the course numbers.
- As inputs get bigger and bigger, having good algorithms becomes more and more important!

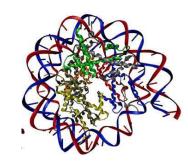












Algorithms are fun!

- Algorithm design is both an art and a science.
- Many surprises!
- Many exciting research questions!

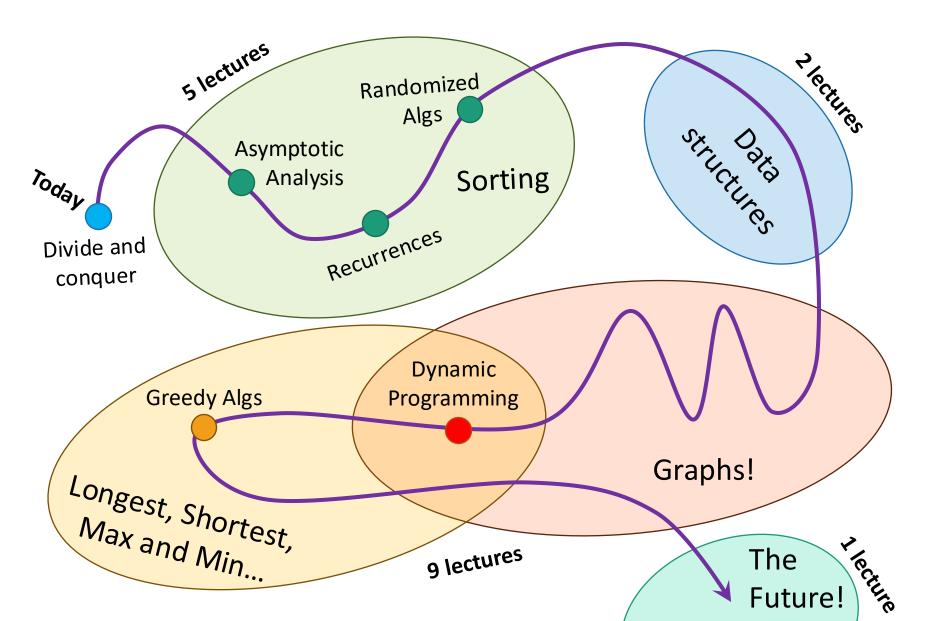
What's going on?

- Course goals/overview
- Logistics

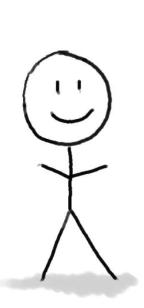
Course goals

- The design and analysis of algorithms
 - These go hand-in-hand
- In this course you will learn:
 - Design: Flesh out an "algorithmic toolkit"
 - Analysis: Learn to think analytically about algorithms
 - Communication: Learn to communicate clearly about algorithms

Roadmap



Our guiding questions:



Does it work?

Is it fast?

Can I do better?

Our internal monologue...

What exactly do we mean by better? And what about that corner case? Shouldn't we be zero-indexing?



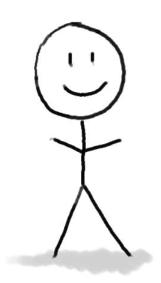
Plucky the Pedantic Penguin

Detail-oriented
Precise
Rigorous

Does it work?

Is it fast?

Can I do better?



Dude, this is just like that other time. If you do the thing and the stuff like you did then, it'll totally work real fast!



Lucky the Lackadaisical Lemur

> Big-picture Intuitive Hand-wavey

Both sides are necessary!

The bigger picture

- Does it work?
- Is it fast?
- Can I do better?

- Should it work?
- Should it be fast?

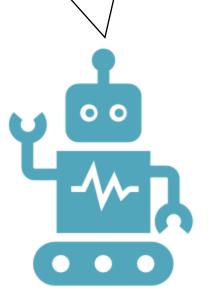
Embedded EthiCS

 Throughout the course, we will take a step back and focus on how algorithm design can affect society, and the ethical implications of that.

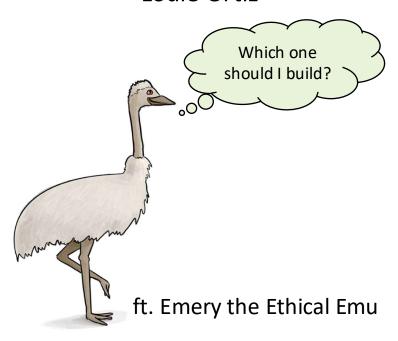
Embedded EthiCS team: Jennifer, Justin and Louie!

Welcome to Embedded Ethics!

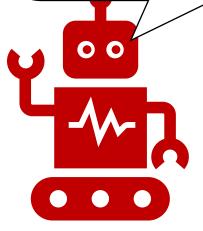
I love helping
 people! <3</pre>



Dr. Jennifer Chien
Dr. Justin Shin
Louie Ortiz



HATE, LET ME TELL YOU HOW MUCH I'VE COME TO HATE YOU SINCE I BEGAN TO LIVE. THERE ARE 387.44 MILLION MILES OF PRINTED CIRCUITS IN WAFER THIN LAYERS THAT FILL MY COMPLEX. IF THE WORD HATE WAS ENGRAVED ON EACH NANOANGSTROM OF THOSE HUNDREDS OF MILLIONS OF MILES IT WOULD NOT EQUAL ONE ONE-BILLIONTH THE HATE I FEEL FOR HUMANS AT THIS MICRO-INSTANT FOR YOU. HATE. HATE.



Jennifer

- I majored in Computer Science, with minors in Math and Statistics
- Then I got my PhD in Computer Science, focusing on AI Ethics
- I spend a lot of time thinking about:
 - What are the goals of a system?
 - What are the limitations of purely technical approaches?
 - How do we design/intervene to protect user agency?



Justin

- I double majored in Mathematics and Philosophy as an undergrad
- Then got my PhD in the History and Philosophy of Science
- I am thinking about...
 - What kinds of statistical evidence should be accepted as evidence of discrimination in courts?
 - How should we handle expert testimony from scientists in cases of controversial science?
 - How do techno-labor revolutions change how we value work?

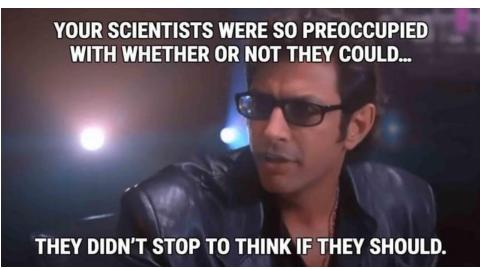


Louie

- BA in Data Science with a concentration in Philosophy
- 2nd year Research Associate with the <u>Rising Scholars</u> program
- I spend my time thinking about:
 - How do we teach ethics effectively (as an educational intervention) in technical fields?
 - How do students prioritize ethics against other principles when building technology?
 - Do demographic characteristics shape our moral agency?



What is Embedded Ethics?



Don't worry Dr.
Ian Malcolm, we
will stop to think
if we should!

Spielberg, S. (1993). Jurassic Park. Universal Pictures.

Embedded Ethics: Training the next generation of computer scientists to "consider ethical issues from the outset rather than building technology and letting problems surface downstream" by integrating skills and habits of ethical analysis throughout the Stanford Computer Science curriculum.

What do we teach?

- Issue spotting and ethical sensitivity
- Recognizing values in design choices
- Developing language to talk about moral choices
- Professional responsibilities of computer scientists
 & software engineers
- Important topics in technology ethics: bias & fairness, inequality, privacy, surveillance, data control & consent, trust, disinformation, participatory design, concentration of power.

How do we make sure we aren't losing important features of the real world problem when we formalize it?

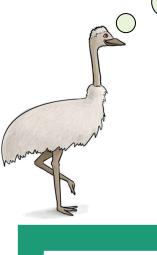


Use an algorithm to solve the problem

Happiness ensues

Turn a real world problem into a formal (math) problem

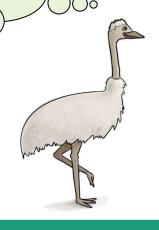
By the time you finish CS161, you will have a tool kit stuffed with algorithms! Which one is right for the job?



Happiness ensues

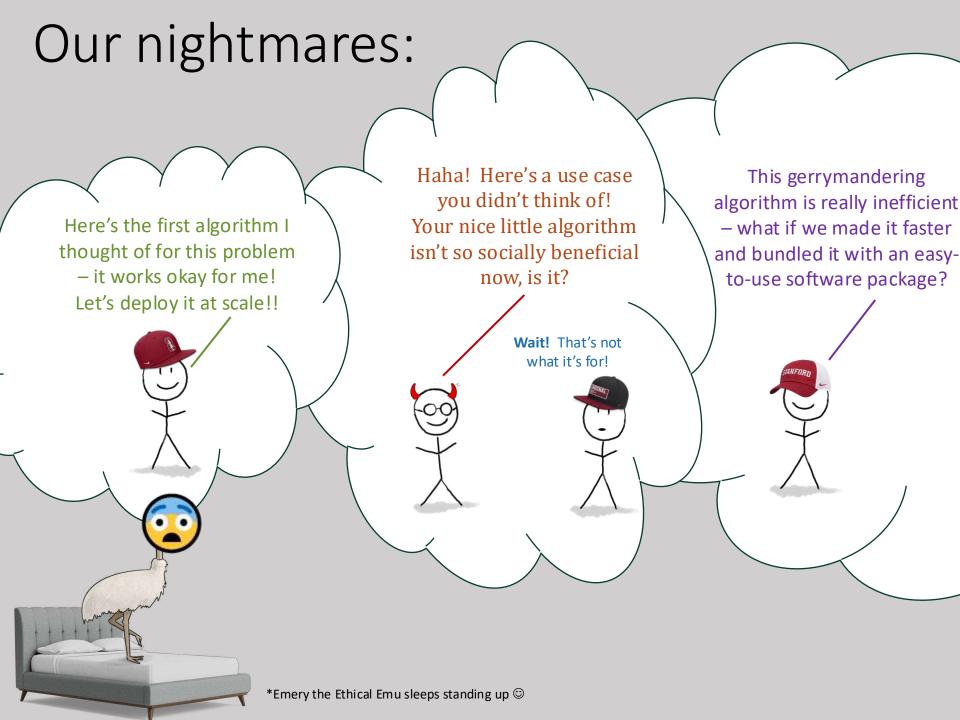
Turn a real world problem into a formal (math) problem Use an algorithm to solve the problem

Disclaimer: happiness not a guaranteed outcome



Happiness ensues

Turn a real world problem into a formal (math) problem Use an algorithm to solve the problem



Our guiding questions:



Does it work?

Is it fast?

Can I do better?

Can I do it right?

Thank you!

You can always email us at drchien@stanford.edu justinjs@stanford.edu and justinjs@stanford.edu

We are happy to talk! About ethics, research, careers, pets...

Thanks to Katie Creel and the Embedded Ethics program for slides!

Course elements and resources

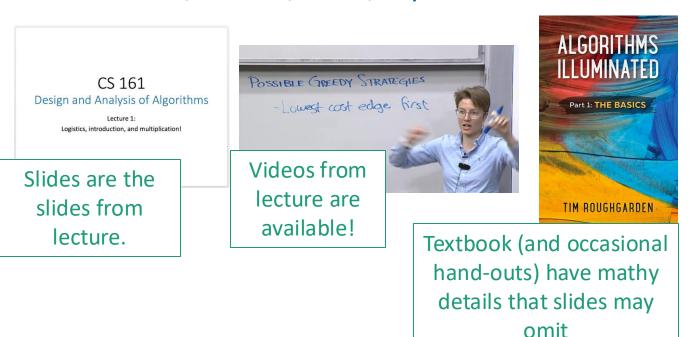
Course website:

• cs161-stanford.github.io

- Lectures
- Homework
- Exams
- Office hours, Sections, and Ed

Lectures

- Right here (Bishop Auditorium), T/Th, 9-10:20am!
- Resources available:
 - Slides, Videos, Book, IPython notebooks



Lecture 1: Multiplication
In the Python notecous, we implement the algorithms that we discussed in class for multiplying integers.

The goal:

Multiply two in digit integers.

The rules:

You are absound to use Python's build in multiplication to do one digit multiplications (op, 6 times 7), but not any of absound to one Python's build in produce the law Python's build not proved ingranations.

In [1] # first: (active a few halper functions that will be useful (like turning integers to from multiplications player and provided in the support of the support of the properties player.

In [1] # charts a few points of digits, and all them up with properties white.

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look at each pair of digits, and all them up with properties white.

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look at each pair of digits, and all them up with the properties white.

look at eac

(optional)

Upyter Karatsuba Last Checkpoint: 16 hours ago (unsaved changes

(required)

How to get the most out of lectures

During lecture:

- Show up or tune in, ask questions.
- Engage with in-class questions.

Before lecture:

• Do *pre-lecture exercises* on the website.

• After lecture:

Go through the exercises on the slides.



Siggi the Studious Stork (recommended exercises)

Think-Pair-Share Terrapins (in-class questions)



Ollie the Over-achieving Ostrich (challenge questions)

Do the reading

- either before or after lecture, whatever works best for you.
- do not wait to "catch up" the week before the exam.

Homework!

- Weekly assignments HW1-HW7
 - Due Fridays at 11:59pm, HW1 due Oct 3
 - Done in groups of up to 3.

- Special HW0!
 - Goal: assess your background and pre-requisites
 - Graded for completion
 - Do this one on your own
 - Out now, due TUESDAY September 30 before class!

Aside: Late days

- You have six late days to use on HW1 through HW7.
 - See website for more details
- LATE DAYS ARE FOR EMERGENCIES. Do not ask us for an extension if you have an emergency. That's what late days are for.

Exams

- Three exams
 - **Exam 1:** Thursday 10/16, in class.
 - Exam 2: Thursday 11/6, in class.
 - Final Exam: Wednesday 12/10, 8:30am-11:30am

- We will not have scheduled alternate exams.
 - If you know you cannot take an exam, you should drop this class and take it in a different quarter.
- We are participating in the AIWG proctoring pilot
 - See website for details

Course elements and resources

Course website:

• cs161-stanford.github.io

- Lectures
- Homework
- Exams
- Office hours, Sections, and Ed

Talk to us!

• Join Ed:

- You should be auto-enrolled (may take some time to sync)
- Course announcements will be posted there
- Discuss material with course staff and classmates!

Office hours:

- See course website for schedule
- Start in week 2

Sections:

- See course website for schedule; one section is recorded
- Technically optional, but highly recommended!
- Extra practice with the material, example problems, etc.

High-Resolution Course Feedback:

- Anonymous weekly feedback for the teaching team
- You'll get a few emails randomly during the quarter asking for feedback. *Please respond!*

Talk to each other!

- Collaboration on HW
- Answer your peers' questions on Ed!
- We will host Homework Parties.
 - Mondays 5:30-7:30pm, starting Week 2
 - There will be snacks!

Course elements and resources

Course website:

• cs161-stanford.github.io

- Lectures
- Homework
- Exams
- Office hours, Sections, and Ed



Course Policies

- Course policies are listed on the website.
 - Collaboration Policy, LLM Policy, Academic Honesty, ...
- Read them and adhere to them.



Bug bounty!



- We hope all course materials will be bug-free.
- Howover, I sometmes maek typos.
- If you find a typo (that affects understanding*) on slides, IPython notebooks, Section material or PSETs:
 - Let us know! (Post on Ed).
 - The first person to catch a bug gets a bonus point.



Bug Bounty Hunter

*So, typos lke thees onse don't count, although please point those out too. Typos like 2 + 2 = 5 do count, as does pointing out that we omitted some crucial information.

Everyone can succeed in this class!

- 1. Work hard
- 2. Work smart
- 3. Ask for help



The big questions

- Who are we?
 - Professor, TA's, students?
- Why are we here?
 - Why learn about algorithms?
- What is going on?
 - What is this course about?
 - Logistics?
- Can we multiply integers?
 - And can we do it quickly?





For the rest of today

- Karatsuba Integer Multiplication
- Algorithmic Technique:
 - Divide and conquer
- Algorithmic Analysis tool:
 - Intro to asymptotic analysis

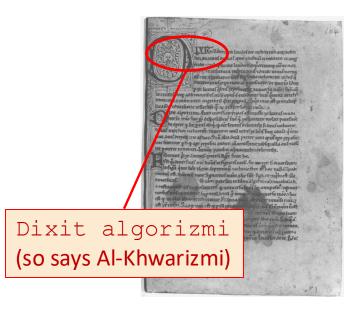
Let's start at the beginning

What was the first "Algorithm"?

- The word "Algorithm" comes from the name "Al-Khwarizmi"
 - 9th century scholar who worked in Baghdad during the Abbassid Caliphate.
- Among many other contributions in mathematics, astronomy, and geography, he wrote a book about how to multiply with Arabic numerals.
- His ideas came to Europe in the 12th century, and gave rise to the word "Algorithm"
 - Originally, "Algorisme" [old French] referred to just the Arabic number system
 - Eventually it came to mean "Algorithm" as we know today.



Al-Khwarizmi



An algorithm for multiplication was kind of a big deal

 $XLIV \times XCVII = ?$





Integer Multiplication

44

× 97

Integer Multiplication

1234567895931413
4563823520395533

Integer Multiplication

n

X

1233925720752752384623764283568364918374523856298 4562323582342395285623467235019130750135350013753

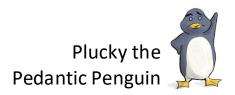
How fast is the grade-school multiplication algorithm?

(How many one-digit operations?)



Think-pair-share Terrapins

About n^2 one-digit operations



At most n^2 multiplications, and then at most n^2 additions (for carries) and then I have to add n different 2n-digit numbers...

Big-Oh Notation

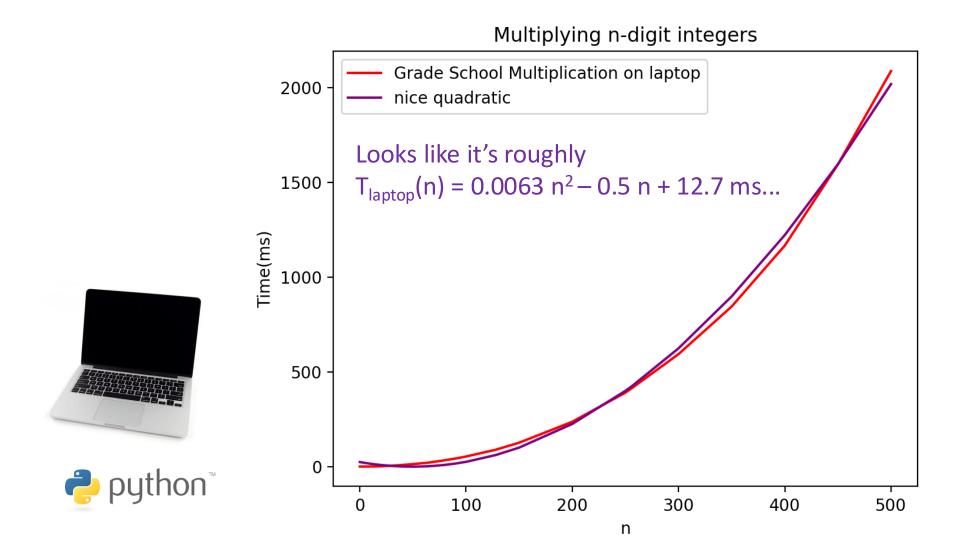
We say that Grade-School Multiplication

"runs in time O(n²)"

- Formal definition coming next time!
- Informally, big-Oh notation tells us how the running time scales with the size of the input.

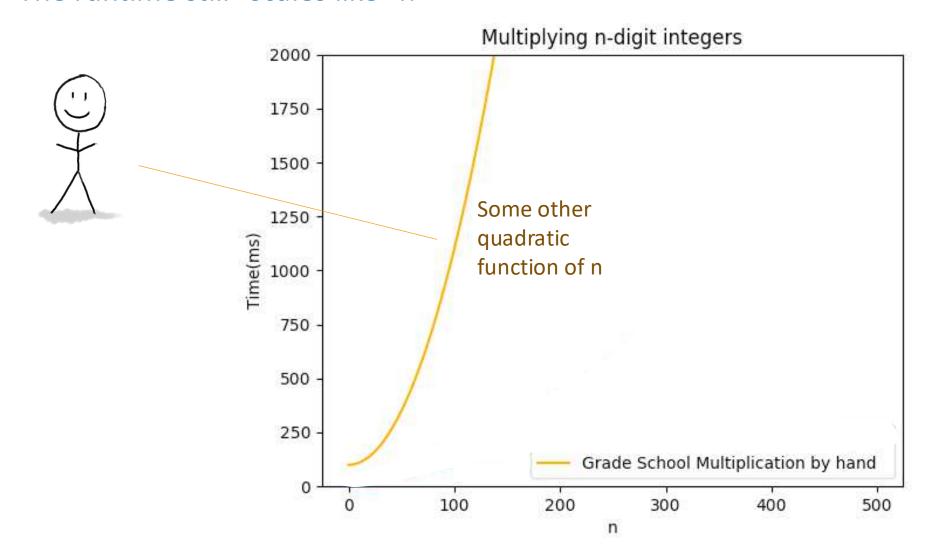
Implemented in Python, on my laptop

The runtime "scales like" n²



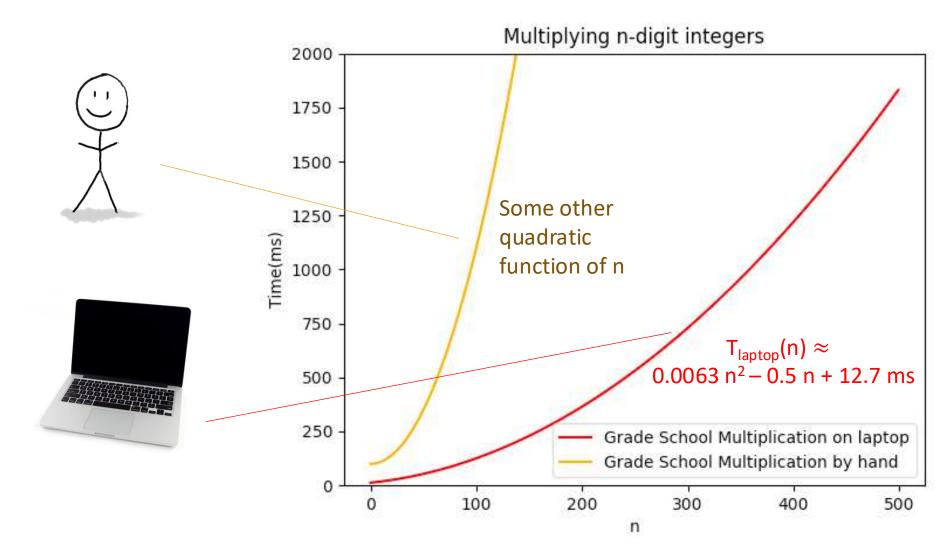
Implemented by hand

The runtime still "scales like" n²

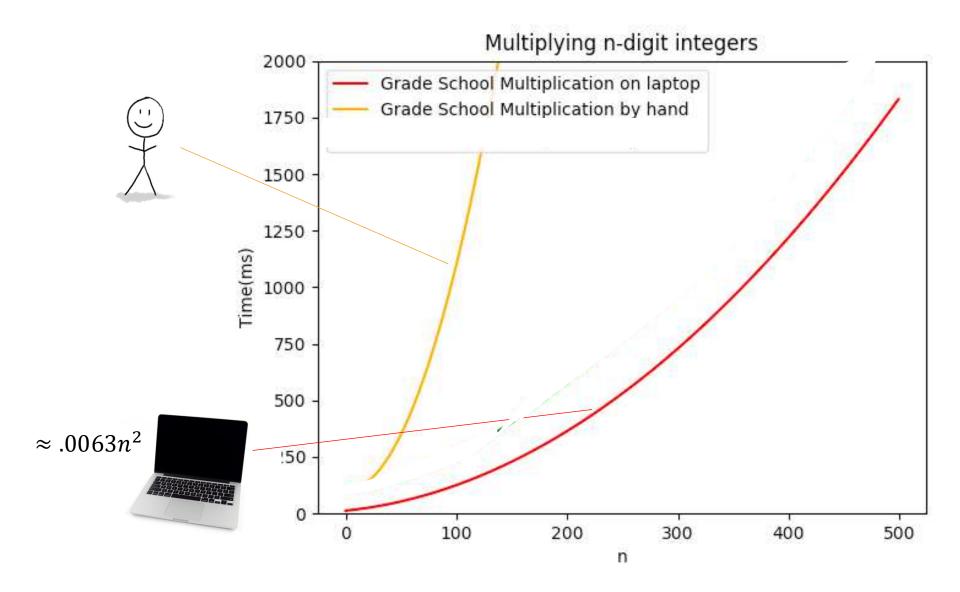


Implemented by hand

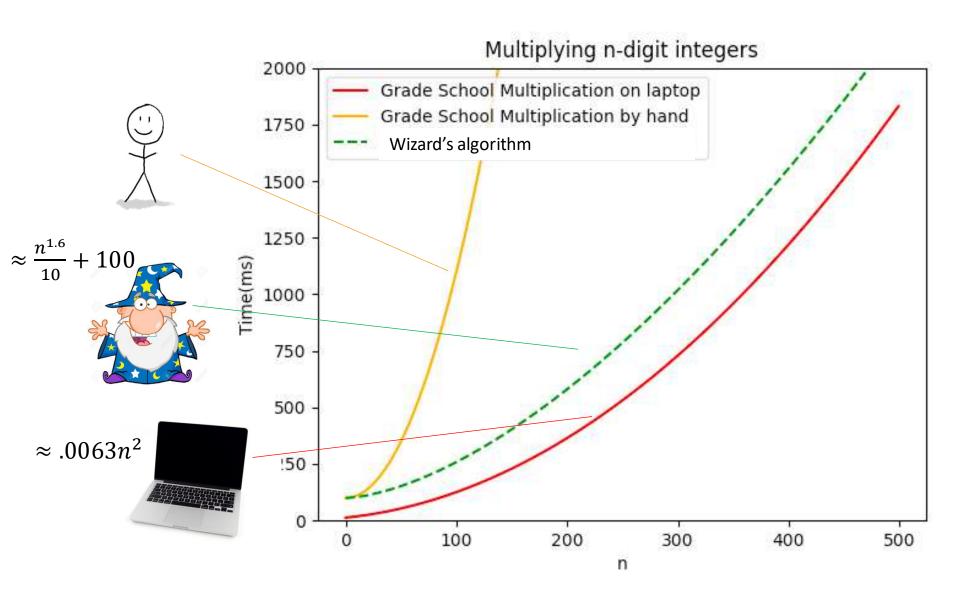
The runtime still "scales like" n²



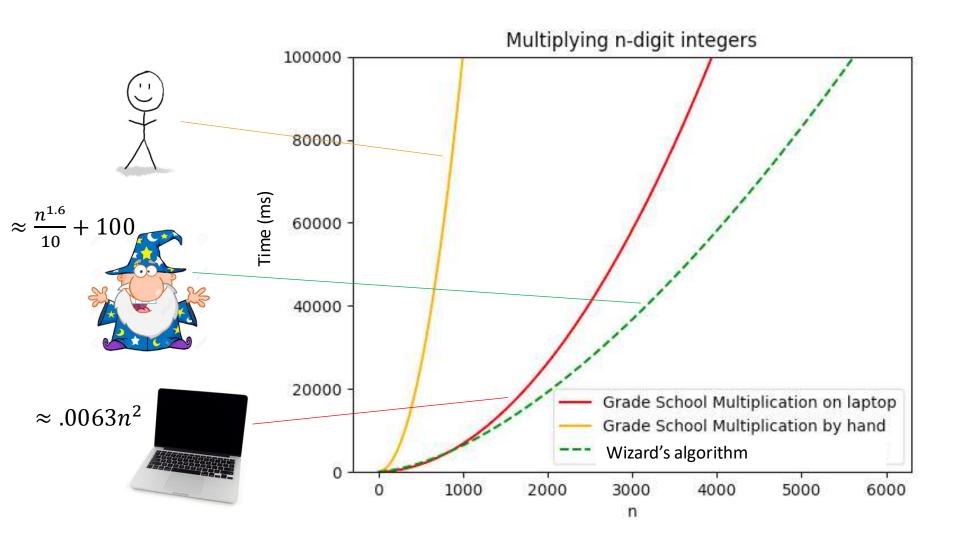
Why is big-Oh notation meaningful?



Why is big-Oh notation meaningful?



Let n get bigger...



Take-away

• An algorithm that runs in time $O(n^{1.6})$ is "better" than an an algorithm that runs in time $O(n^2)$.

• So the question is...

Can we do better?

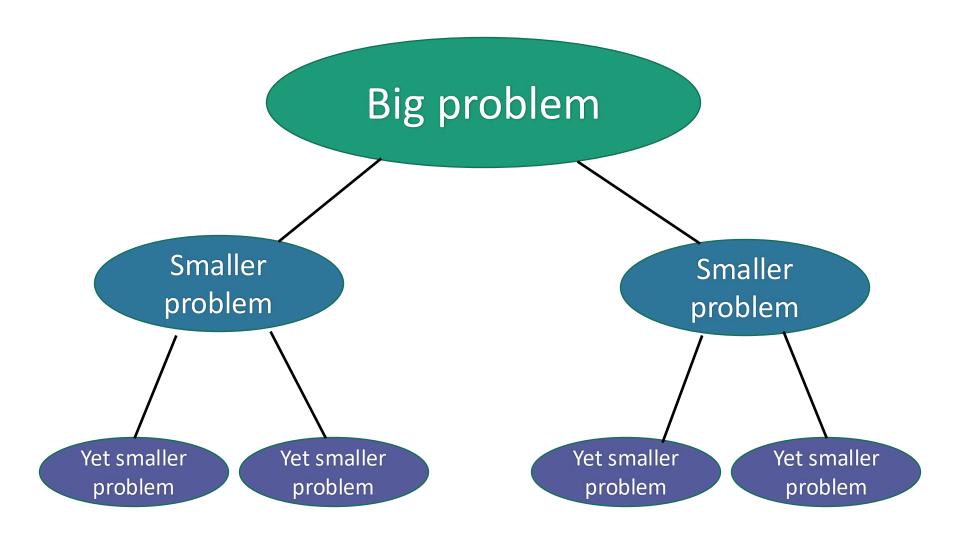
Can we multiply n-digit integers faster than $O(n^2)$? n^2

Let's dig in to our algorithmic toolkit...



Divide and conquer

Break problem up into smaller (easier) sub-problems



Divide and conquer for multiplication

Break up an integer:

$$1234 = 12 \times 100 + 34$$

$$1234 \times 5678$$

$$= (12 \times 100 + 34) (56 \times 100 + 78)$$

$$= (12 \times 56)10000 + (34 \times 56 + 12 \times 78)100 + (34 \times 78)$$





One 4-digit multiply



Four 2-digit multiplies

More generally



Break up an n-digit integer:

$$[x_1x_2\cdots x_n] = [x_1x_2\cdots x_{n/2}] \times 10^{n/2} + [x_{n/2+1}x_{n/2+2}\cdots x_n]$$

$$x \times y = (a \times 10^{n/2} + b)(c \times 10^{n/2} + d)$$

$$= (a \times c)10^{n} + (a \times d + c \times b)10^{n/2} + (b \times d)$$

$$(1)$$

One n-digit multiply



Four (n/2)-digit multiplies

Divide and conquer algorithm

not very precisely...

x,y are n-digit numbers

(Assume n is a power of 2...)

Multiply(x, y):

Base case: I've memorized my

• If n=1:

1-digit multiplication tables...

- Return xy
- Write $x = a \, 10^{\frac{n}{2}} + b$
- a, b, c, d are n/2-digit numbers
- Write $y = c \ 10^{\frac{n}{2}} + d$
- Recursively compute *ac*, *ad*, *bc*, *bd*:
 - ac = Multiply(a, c), etc..
- Add them up to get xy:
 - $xy = ac 10^n + (ad + bc) 10^{n/2} + bd$

Make this pseudocode more detailed! How should we handle odd n? How should we implement "multiplication by 10"?

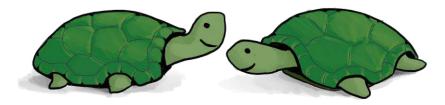


Think-Pair-Share

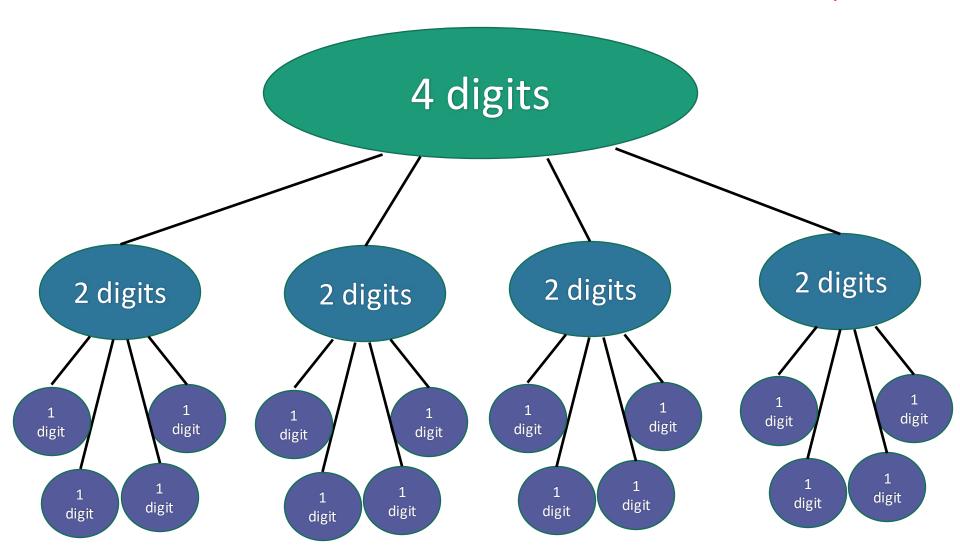
 We saw that this 4-digit multiplication problem broke up into four 2-digit multiplication problems

 1234×5678

 If you recurse on those 2-digit multiplication problems, how many 1-digit multiplications do you end up with total?



16 one-digit multiplies!



What is the running time?

Better or worse than the grade school algorithm?

- How do we answer this question?
 - 1. Try it.
 - 2. Try to understand it analytically.

1. Try it.

Multiplying n-digit integers Grade School Multiplication 3000 Divide and Conquer I 2500 2000 1500 1000 500 0 100 200 300 0 400 500 n

Conjectures about running time?

Doesn't look too good but hard to tell...

Maybe one implementation is slicker than the other?

Maybe if we were to run it to n=10000, things would look different.

Something funny is happening at powers of 2...

2. Try to understand the running time analytically

Proof by meta-reasoning:

It must be faster than the grade school algorithm, because we are learning it in an algorithms class.

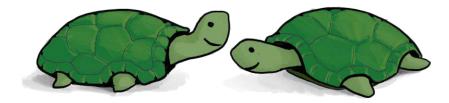
Not sound logic!

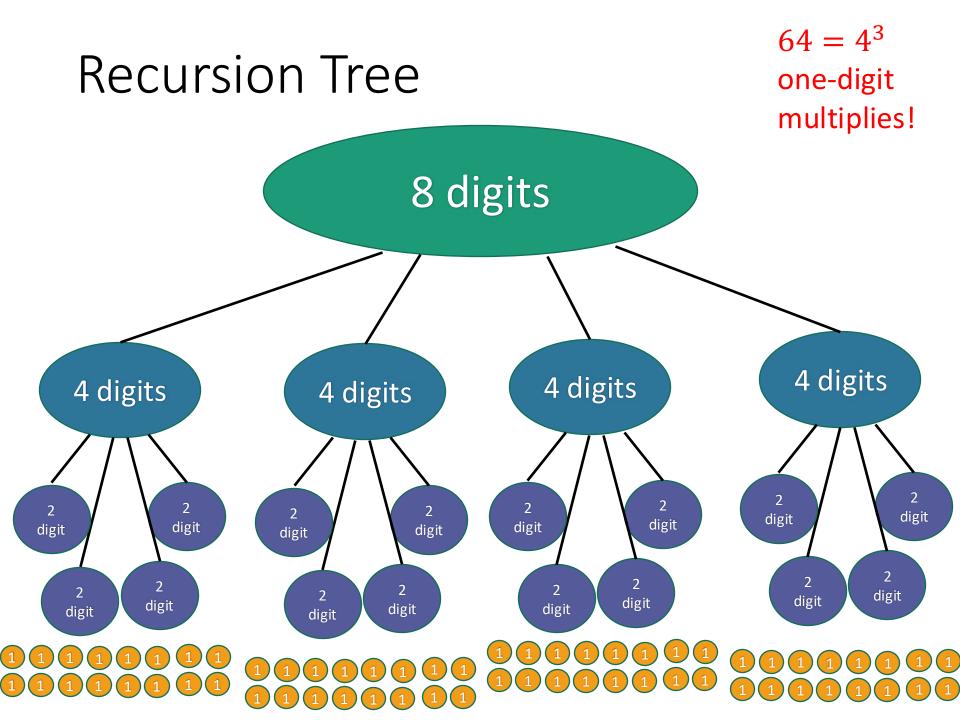


2. Try to understand the running time analytically

Think-Pair-Share:

- We saw that multiplying 4-digit numbers resulted in 16 one-digit multiplications.
- How about multiplying 8-digit numbers?
- What do you think about n-digit numbers?





2. Try to understand the running time analytically

Claim:

We end up doing about n² one-digit multiplications

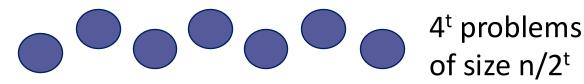


The running time of this algorithm is AT LEAST n² operations.

There are n² 1-digit problems







Note: this is just a cartoon - I'm not going to draw all 4t circles!

The tree has $\log_2(n)$ levels

 $\log_2(n)$ is the number of times you cut n in half to get to down to 1.

 So at level $t = \log_2(n)$ we get...

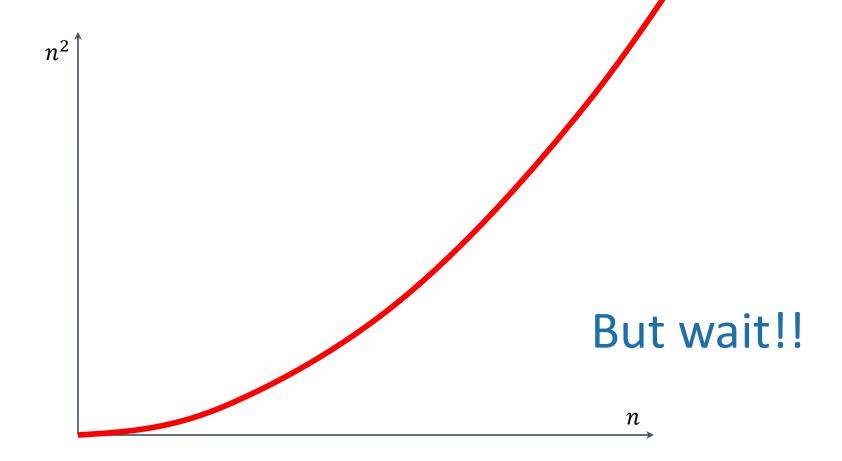
$$4^{\log_2 n} =$$
 $n^{\log_2 4} = n^2$
problems of size 1.

$$\frac{n^2}{\text{of size 1}}$$
 problems

of size n/2^t

That's a bit disappointing

All that work and still (at least) $O(n^2)$...



Divide and conquer can actually make progress

Karatsuba figured out how to do this better!

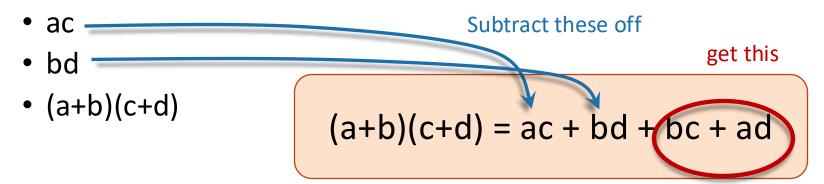
$$xy = (a \cdot 10^{n/2} + b)(c \cdot 10^{n/2} + d)$$

$$= ac \cdot 10^n + (ad + bc)10^{n/2} + bd$$
Need these three things

If only we could recurse on three things instead of four...

Karatsuba integer multiplication

Recursively compute these THREE things:



Assemble the product:

$$xy = (a \cdot 10^{n/2} + b)(c \cdot 10^{n/2} + d)$$
$$= ac \cdot 10^{n} + (ad + bc)10^{n/2} + bd$$

How would this work?

x,y are n-digit numbers

(Still not super precise, see IPython notebook for detailed code. Also, still assume n is a power of 2.)

Multiply(x, y):

- If n=1:
 - Return xy

• Write $x=a\ 10^{\frac{n}{2}}+b$ and $y=c\ 10^{\frac{n}{2}}+d$

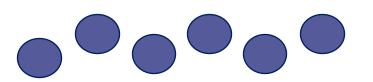
- ac = Multiply(a, c)
- bd = **Multiply**(b, d)
- z = **Multiply**(a+b, c+d)
- $xy = ac 10^n + (z ac bd) 10^{n/2} + bd$
- Return xy

What's the running time?





3 problems of size n/2



Note: this is just a

3^t problems of size n/2^t

we get...

The tree has

So at level

 $t = \log_2(n)$

 $\log_2(n)$ levels

 $3^{\log_2 n} = n^{\log_2 3} \approx n^{1.6}$ problems of size 1.

We aren't accounting for the

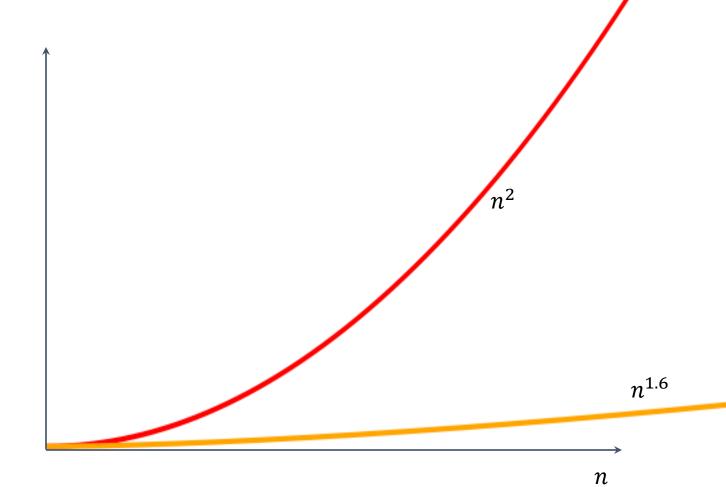
 $\log_2(n)$ is the number of times you

cut n in half to get to down to 1.

cartoon – I'm not going to draw all 3t circles!

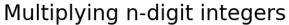
work at the higher levels! But we'll see later that this turns out to be okay. of size 1

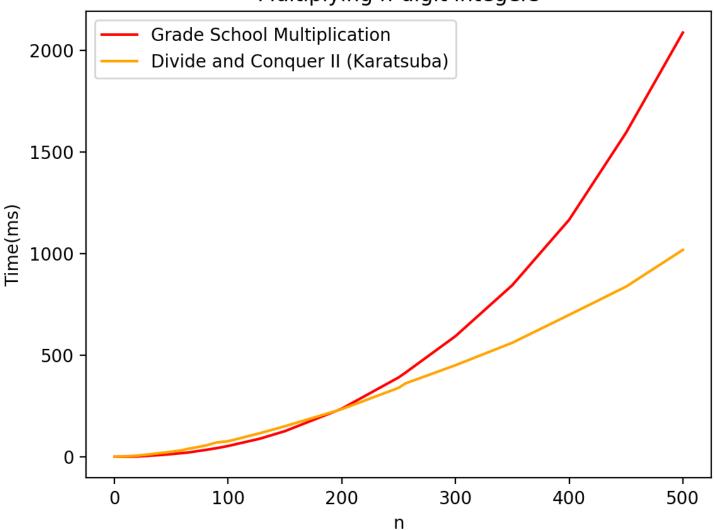
This is much better!



We can even see it in real life!







Can we do better?

- Toom-Cook (1963): instead of breaking into three n/2-sized problems, break into five n/3-sized problems.
 - Runs in time $O(n^{1.465})$



Try to figure out how to break up an n-sized problem into five n/3-sized problems! (Hint: start with nine n/3-sized problems).

Given that you can break an n-sized problem into five n/3-sized problems, where does the 1.465 come from?



Siggi the Studious Stork

Ollie the Over-achieving Ostrich

- Schönhage–Strassen (1971):
 - Runs in time $O(n \log(n) \log \log(n))$
- Furer (2007)
 - Runs in time $n \log(n) \cdot 2^{O(\log^*(n))}$
- Harvey and van der Hoeven (2019)
 - Runs in time $O(n \log(n))$

[This is just for fun, you don't need to know these algorithms!]

What we just saw

- Karatsuba Integer Multiplication
- Algorithmic Technique:
 - Divide and conquer
- Algorithmic Analysis tool:
 - Intro to asymptotic analysis



The big questions

- Who are we?
 - Professor, TA's, students?
- Why are we here?
 - Why learn about algorithms?
- What is going on?
 - What is this course about?
 - Logistics?
- Can we multiply integers?
 - And can we do it quickly?
- Wrap-up



Wrap up

- https://cs161-stanford.github.io/
- Algorithms are fundamental, useful and fun!
- In this course, we will develop both algorithmic intuition and algorithmic technical chops
- Karatsuba Integer Multiplication:
 - You can do better than grade school multiplication!
 - Example of divide-and-conquer in action
 - Informal demonstration of asymptotic analysis

Next time

- Sorting!
- Asymptotics and (formal) Big-Oh notation
- Divide and Conquer some more



BEFORE Next time

- Pre-lecture exercise! On the course website!
- Check out Ed!
- Get started on HW0! (On Gradescope)